

Contrary Pairs in Pseudo-Boolean Algebra

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Abstract.

The lattice operations for split subsets are considered when the values of membership function belong to the pseudo-boolean algebra.

Let us consider a set of pairs of kind $(\mu I_A, (1-\mu)I_A)$, where $A \subseteq \Omega$, Ω is universal domain, A - his fixed subset, μ - any mapping of Ω into $[0,1]$ and I_A - characteristic function or indicator of A /1/.

Introduce \wedge (greatest lower bound) and \vee (least upper bound) operations traditionally componentwise:

$$(X, Y) \vee (Z, T) = (X \vee Z, Y \vee T)$$

$$(X, Y) \wedge (Z, T) = (X \wedge Z, Y \wedge T)$$

Take arbitrarily two pairs $(\mu I_A, (1-\mu)I_A)$, $(\nu I_A, (1-\nu)I_A)$, $\mu, \nu \in [0,1]^\Omega$ and consider their l.u.b. and g.l.b:

$$\begin{aligned} ((\mu I_A \vee \nu I_A)(x), ((1-\mu)I_A \vee (1-\nu)I_A)(x)) &= \\ &= \begin{cases} (0,0) & \text{if } x \notin A \\ (\nu(x), 1-\mu(x)) & \text{if } x \in A, \mu(x) \leq \nu(x) \\ (\mu(x), 1-\nu(x)) & \text{if } x \in A, \mu(x) > \nu(x) \end{cases} \\ ((\mu I_A \wedge \nu I_A)(x), ((1-\mu)I_A \wedge (1-\nu)I_A)(x)) &= \\ &= \begin{cases} (0,0) & \text{if } x \notin A \\ (\mu(x), 1-\nu(x)) & \text{if } x \in A, \mu(x) \leq \nu(x) \\ (\nu(x), 1-\mu(x)) & \text{if } x \in A, \mu(x) > \nu(x) \end{cases} \end{aligned}$$

Thus the set of kind $(\mu I_A, (1-\mu)I_A)$ is not closed with respect to operations \wedge and \vee . If we want to keep closely, then operations will be defined as follows: $(X, Y) \vee (Z, T) = (X \vee Z, Y \wedge T)$

In this case, we have:

$$\begin{aligned} ((\mu I_A \vee \nu I_A)(x), ((1-\mu)I_A \wedge (1-\nu)I_A)(x)) &= \\ &= \begin{cases} (0,0) & \text{if } x \notin A \\ (\nu(x), 1-\nu(x)) & \text{if } x \in A, \mu(x) \leq \nu(x) \\ (\mu(x), 1-\mu(x)) & \text{if } x \in A, \mu(x) > \nu(x) \end{cases} \\ ((\mu I_A \wedge \nu I_A)(x), ((1-\mu)I_A \vee (1-\nu)I_A)(x)) &= \\ &= \begin{cases} (0,0) & \text{if } x \notin A \\ (\mu(x), 1-\mu(x)) & \text{if } x \in A, \mu(x) \leq \nu(x) \\ (\nu(x), 1-\nu(x)) & \text{if } x \in A, \mu(x) > \nu(x) \end{cases} \end{aligned}$$

\wedge and \vee operations induce a following partial ordering relation between the pairs

$$\begin{aligned} (X, Y) \leq (Z, T) &\Leftrightarrow (X, Y) \vee (Z, T) = (Z, T) \Leftrightarrow \\ &\Leftrightarrow (X, Y) \wedge (Z, T) = (X, Y) \Leftrightarrow (X(u) \leq Z(u)) \& (Y(u) \geq T(u)) \\ &X, Y, Z, T \in [0,1]^A, u \in A. \end{aligned}$$

The complement $(X, Y)^c$ of pair (X, Y) is defined in this way: $(X, Y)^c = (I_A \setminus X, I_A \setminus Y)$

Denote by a symbol L the lattice of fuzzy subsets of Ω and by symbols L' the lattice of same subsets with reverse order. It is fairly straitforward to show that the following theorem holds:

THEOREM 1. Pairs of kind $(\mu I_A, (1-\mu)I_A)$ for fixed $A \subseteq \Omega$ and any $\mu \in [0,1]^\Omega$ form complemented distributive lattice, which is a sublattice of $L \times L'$, satisfying de Morgan's law; the complement is involutory and order reversing.

Now suppose, that the values of membership function μ belong to the pseudo-boolean algebra $B = \langle B, \leq \rangle$ /2/. As intuitionistic negation is not, in general, involutory, the unique splitting into contrary pair is impossible /3/. In this case, the construction defined below perhaps proved to be useful.

We introduce the " \approx " relation between elements of B and between their pairs as follows:

DEFINITION 1. $(a \approx b) \stackrel{df}{=} (a^* = b^*)$, $a, b \in B$. Here $()^*$ denotes pseudo-complement.

2. $((a, b) \approx (c, d)) \stackrel{df}{=} ((a \approx c) \text{ and } (b \approx d))$, $a, b, c, d \in B$.

3. $(a, b)^* \stackrel{df}{=} (a^*, b^*)$, $a, b \in B$.

It is evident that " \approx " is an equivalence relation.

THEOREM 2. If $a_1 \approx a_2$, $b_1 \approx b_2$ then
 $((a_1, a_1^*) \vee (b_1, b_1^*)) \approx ((a_2, a_2^*) \vee (b_2, b_2^*))$, $(a_1, a_1^*)^* \approx (a_2, a_2^*)^*$

PROOF.

$$(a_1, a_1^*) \vee (b_1, b_1^*) = (a_1 \vee b_1, a_1^* \wedge b_1^*) = (a_1 \vee b_1, a_2^* \wedge b_2^*) \\ (a_2, a_2^*) \vee (b_2, b_2^*) = (a_2 \vee b_2, a_2^* \wedge b_2^*), \quad (a_1 \vee b_1)^* = a_1^* \wedge b_1^* = a_2^* \wedge b_2^* = (a_2 \vee b_2)^*.$$

2. follows immediately from definitions.

Let $B' \subset B$ be the set of elements satisfying the following condition:

$$(a \wedge b)^* = a^* \vee b^*$$

THEOREM 3. If $a, b, c, d \in B'$, $a_1 \approx a_2$, $b_1 \approx b_2$ then
 $((a_1, a_1^*) \wedge (b_1, b_1^*)) \approx ((a_2, a_2^*) \wedge (b_2, b_2^*))$

Proof by analogy with case 1. of theorem given above. Thus, we can consider the factor set $B' \times B' / \approx$.

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